

CSC 2400: Computer Systems

Assembly: Accessing Memory

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Review:

Instructions to Recognize

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Arithmetic Instructions (1)

- Two Operand Instructions

<code>addl Src, Dest</code>	$Dest = Dest + Src$	
<code>subl Src, Dest</code>	$Dest = Dest - Src$	
<code>imull Src, Dest</code>	$Dest = Dest * Src$	
<code>sall Src, Dest</code>	$Dest = Dest \ll Src$	
<code>sarl Src, Dest</code>	$Dest = Dest \gg Src$	Arithmetic
<code>shrl Src, Dest</code>	$Dest = Dest \gg Src$	Logical
<code>xorl Src, Dest</code>	$Dest = Dest \wedge Src$	
<code>andl Src, Dest</code>	$Dest = Dest \& Src$	
<code>orl Src, Dest</code>	$Dest = Dest Src$	

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Arithmetic Instructions (2)

- One Operand Instructions

<code>incl Dest</code>	$Dest = Dest + 1$
<code>decl Dest</code>	$Dest = Dest - 1$
<code>negl Dest</code>	$Dest = -Dest$
<code>notl Dest</code>	$Dest = \sim Dest$

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Compare and Test Instructions

`cmpl` *Src, Dest*
Compute *Dest-Src* without setting *Dest*

`testl` *Src, Dest*
Compute *Dest&Src* without setting *Dest*

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Jump Instructions

- Jump depending on the result of the previous arithmetic instruction:

Jump	Description
<code>jmp</code>	Unconditional
<code>je</code>	Equal / Zero
<code>jne</code>	Not Equal / Not Zero
<code>js</code>	Negative
<code>jns</code>	Nonnegative
<code>jg</code>	Greater (Signed)
<code>jge</code>	Greater or Equal (Signed)
<code>jl</code>	Less (Signed)
<code>jle</code>	Less or Equal (Signed)
<code>ja</code>	Above (unsigned)
<code>jb</code>	Below (unsigned)

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Accessing Memory

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Lecture Goals

- Help you learn...
 - To manipulate data of various sizes
 - To leverage more sophisticated addressing modes
- Focusing on the assembly-language code
 - Rather than the layout of memory for storing data
- Why?
 - Understand the relationship to data types and common programming constructs in higher-level languages

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Variable Sizes in High-Level Language

- C data types vary in size
 - Character: 1 byte
 - Short, int, and long: varies, depending on the computer
 - Float and double: varies, depending on the computer
 - Pointers: typically 4 bytes
- Programmer-created types
 - Struct: arbitrary size, depending on the fields
- Arrays
 - Multiple consecutive elements of some fixed size
 - Where each element could be a struct

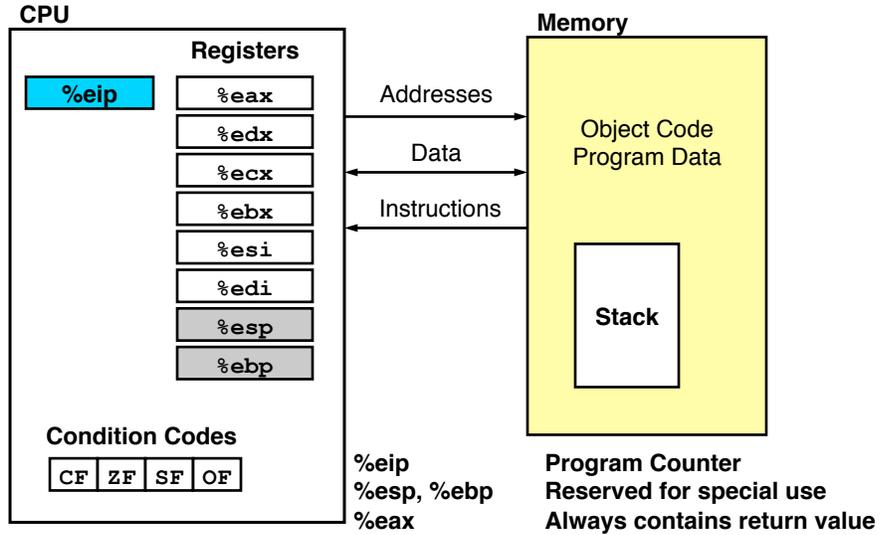
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Supporting Different Sizes in IA-32

- Three main data sizes
 - Byte (b): 1 byte
 - Word (w): 2 bytes
 - Long (l): 4 bytes
- Separate assembly-language instructions
 - E.g., `addb`, `addw`, and `addl`
- Separate ways to access (parts of) a register
 - E.g., `%ah` or `%al`, `%ax`, and `%eax`
- Larger sizes (e.g., struct)
 - Manipulated in smaller byte, word, or long units

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IA-32 Architecture



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IA-32 General Purpose Registers

31	15	8	7	0	16-bit	32-bit
	AH	AL			AX	EAX
	BH	BL			BX	EBX
	CH	CL			CX	ECX
	DH	DL			DX	EDX
	SI					ESI
	DI					EDI

General-purpose registers

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Byte Order in Multi-Byte Entities

- Intel is a **little endian** architecture
 - Least significant byte of multi-byte entity is stored at lowest memory address
 - “Little end goes first”

The int 5 at address 1000:

1000	00000101
1001	00000000
1002	00000000
1003	00000000

- Some other systems use **big endian**
 - Most significant byte of multi-byte entity is stored at lowest memory address
 - “Big end goes first”

The int 5 at address 1000:

1000	00000000
1001	00000000
1002	00000000
1003	00000101

Little Endian Example

```
int main(void) {
    int i=0x003377ff, j;
    unsigned char *p = (unsigned char *) &i;
    for (j=0; j<4; j++)
        printf("Byte %d: %x\n", j, p[j]);
}
```

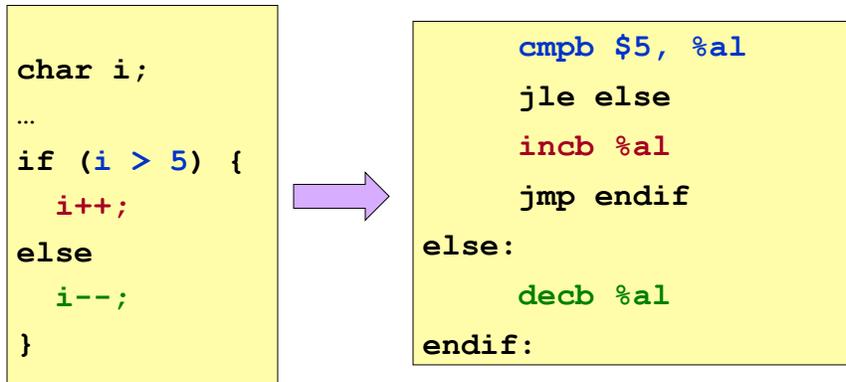
Output on a
little-endian
machine

Byte 0: ff
Byte 1: 77
Byte 2: 33
Byte 3: 0

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C Example: One-Byte Data

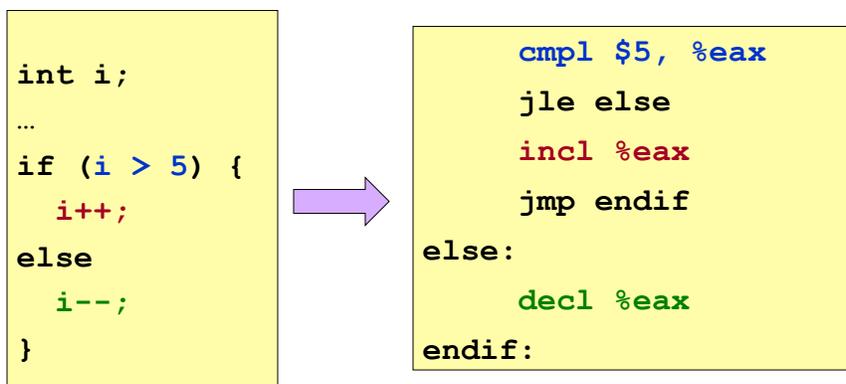
Global *char* variable *i* is in *%al*,
the *lower byte* of the “A” register.



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C Example: Four-Byte Data

Global *int* variable *i* is in *%eax*,
the *full 32 bits* of the “A” register.



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Loading and Storing Data

- Processors have many ways to access data
 - Known as “addressing modes”
 - Two simple ways seen in previous examples
- Immediate addressing
 - Example: `movl $0, %ecx`
 - Data (e.g., number “0”) embedded in the instruction
 - Initialize register ECX with zero
- Register addressing
 - Example: `movl %edx, %ecx`
 - Choice of register(s) embedded in the instruction
 - Copy value in register EDX into register ECX

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Direct Addressing

- Load or store from a particular memory location
 - Memory address is embedded in the instruction
 - Instruction reads from or writes to that address
- IA-32 example: `movl 2000, %ecx`
 - Four-byte variable located at address 2000
 - Read four bytes starting at address 2000
 - Load the value into the ECX register
- Useful when the address is known in advance
 - Global variables in the Data or BSS sections
- Can use a label for (human) readability
 - E.g., “i” to allow “`movl i, %eax`”

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Indirect Addressing

- Load or store from a previously-computed address
 - Register with the address is embedded in the instruction
 - Instruction reads from or writes to that address
- IA-32 example: `movl (%eax), %ecx`
 - EAX register stores a 32-bit address (e.g., 2000)
 - Read long-word variable stored at that address
 - Load the value into the ECX register
- Useful when address is not known in advance
 - Dynamically allocated data referenced by a pointer
 - The “(%eax)” essentially dereferences a pointer

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Base Pointer Addressing

- Load or store with an offset from a base address
 - Register storing the base address
 - Fixed offset also embedded in the instruction
 - Instruction computes the address and does access
- IA-32 example: `movl 8(%eax), %ecx`
 - EAX register stores a 32-bit base address (e.g., 2000)
 - Offset of 8 is added to compute address (e.g., 2008)
 - Read long-word variable stored at that address
 - Load the value into the ECX register
- Useful when accessing part of a larger variable
 - Specific field within a “struct”
 - E.g., if “age” starts at the 8th byte of “student” record

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Indexed Addressing

- Load or store with an offset and multiplier
 - Fixed based address embedded in the instruction
 - Offset computed by multiplying register with constant
 - Instruction computes the address and does access
- IA-32 example: `movl 2000(,%eax,4), %ecx`
 - Index register EAX (say, with value of 10)
 - Multiplied by a multiplier of 1, 2, 4, or 8 (say, 4)
 - Added to a fixed base of 2000 (say, to get 2040)
- Useful to iterate through an array (e.g., `a[i]`)
 - Base is the start of the array (i.e., “a”)
 - Register is the index (i.e., “i”)
 - Multiplier is the size of the element (e.g., 4 for “int”)

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Indexed Addressing Example

```
int a[20]; ← global variable
```

```
int i, sum=0;  
for (i=0; i<20; i++)  
    sum += a[i];
```

EAX: i
EBX: sum
ECX: temporary

```
movl $0, %eax  
movl $0, %ebx  
sumloop:  
    movl a(,%eax,4), %ecx  
    addl %ecx, %ebx  
    incl %eax  
    cmpl $19, %eax  
    jle sumloop
```

Indexed Address: D(Base,Index,Scale)

$$\text{Address} = \text{Base} + \text{Index} * \text{Scale} + \text{D}$$

Examples:

- Displacement `movl foo, %ebx`
- Base `movl (%eax), %ebx`
- Base + disp `movl foo(%eax), %ebx`
`movl 1(%eax), %ebx`
- (Index * scale) + disp `movl (,%eax,4), %ebx`
- Base + (index * scale) + disp `movl foo(%edx,%eax,4), %ebx`

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Address Computation Examples

<code>%edx</code>	<code>0xf000</code>
<code>%ecx</code>	<code>0x100</code>

Expression	Computation	Address
<code>0x8(%edx)</code>	<code>0xf000 + 0x8</code>	<code>0xf008</code>
<code>(%edx,%ecx)</code>	<code>0xf000 + 0x100</code>	<code>0xf100</code>
<code>(%edx,%ecx,4)</code>	<code>0xf000 + 4*0x100</code>	<code>0xf400</code>
<code>0x80(,%edx,2)</code>	<code>2*0xf000 + 0x80</code>	<code>0x1e080</code>

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Address Computation Instruction

- `leal Src, Dest`
 - `Src` is address mode expression
 - Set `Dest` to address denoted by expression
- Uses
 - Computing address without doing memory reference
 - E.g., translation of `p = &x[i];`
 - Computing arithmetic expressions of the form $x + k*y + z$
 - $k = 1, 2, 4, \text{ or } 8.$
 - z is an 8-bit signed constant

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Using `leal` for Arithmetic Expressions

```
int arith
(int x, int y, int z)
{
    int t1 = x+y;
    int t2 = z+t1;
    int t3 = x+4;
    int t4 = y * 48;
    int t5 = t3 + t4;
    int rval = t2 * t5;
    return rval;
}
```

```
arith:
    pushl %ebp
    movl %esp,%ebp
} Set Up

    movl 8(%ebp),%eax
    movl 12(%ebp),%edx
    leal (%edx,%eax),%ecx
    leal (%edx,%edx,2),%edx
    sall $4,%edx
    addl 16(%ebp),%ecx
    leal 4(%edx,%eax),%eax
    imull %ecx,%eax
} Body

    movl %ebp,%esp
    popl %ebp
    ret
} Finish
```

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Understanding arith

```
int arith
(int x, int y, int z)
{
  int t1 = x+y;
  int t2 = z+t1;
  int t3 = x+4;
  int t4 = y * 48;
  int t5 = t3 + t4;
  int rval = t2 * t5;
  return rval;
}
```

x is at address ebp+8
y is at address ebp+12
x is at address ebp+16

To be explained in next lecture

```
movl 8(%ebp),%eax      # eax =
movl 12(%ebp),%edx     # edx =
leal (%edx,%eax),%ecx  # ecx =
leal (%edx,%edx,2),%edx # edx =
sall $4,%edx          # edx =
addl 16(%ebp),%ecx     # ecx =
leal 4(%edx,%eax),%eax # eax =
imull %ecx,%eax       # eax =
```

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Understanding arith

```
int arith
(int x, int y, int z)
{
  int t1 = x+y;
  int t2 = z+t1;
  int t3 = x+4;
  int t4 = y * 48;
  int t5 = t3 + t4;
  int rval = t2 * t5;
  return rval;
}
```

x is at address ebp+8
y is at address ebp+12
x is at address ebp+16

To be explained in next lecture

```
movl 8(%ebp),%eax      # eax = x
movl 12(%ebp),%edx     # edx = y
leal (%edx,%eax),%ecx  # ecx = x+y (t1)
leal (%edx,%edx,2),%edx # edx = 3*y
sall $4,%edx          # edx = 48*y (t4)
addl 16(%ebp),%ecx     # ecx = z+t1 (t2)
leal 4(%edx,%eax),%eax # eax = 4+t4+x (t5)
imull %ecx,%eax       # eax = t5*t2 (return value) 29
```

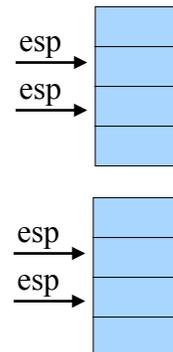
Data Access Methods: Summary

- **Immediate addressing:** data stored in the instruction itself
 - `movl $10, %ecx`
- **Register addressing:** data stored in a register
 - `movl %eax, %ecx`
- **Direct addressing:** address stored in instruction
 - `movl foo, %ecx`
- **Indirect addressing:** address stored in a register
 - `movl (%eax), %ecx`
- **Base pointer addressing:** includes an offset as well
 - `movl 4(%eax), %ecx`
- **Indexed addressing:** instruction contains base address, and specifies an index register and a multiplier (1, 2, 4, or 8)
 - `movl 2000(%eax,1), %ecx`

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Data Transfer Instructions

- **`mov{b,w,l} source, dest`**
 - General move instruction
- **`push{w,l} source`**
`pushl %ebx` # equivalent instructions
`subl $4, %esp`
`movl %ebx, (%esp)`
- **`pop{w,l} dest`**
`popl %ebx` # equivalent instructions
`movl (%esp), %ebx`
`addl $4, %esp`



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Conclusions

- Accessing data
 - Byte, word, and long-word data types
 - Wide variety of addressing modes
- Next time
 - Calling functions, using the stack